

Written Exam in  
**Data compression**  
**TSBK08**

27th August 2022 8:00 - 12:00

<b>Location:</b>	TER2
<b>Examiner:</b>	Harald Nautsch
<b>Teacher:</b>	Harald Nautsch, 0701718715
<b>Department:</b>	ISY
<b>Module:</b>	TEN1
<b>Number of problems:</b>	6
<b>Number of pages:</b>	4
<b>Permitted equipment:</b>	Calculator, general English dictionaries
<b>Other:</b>	Answers can be given in English or in Swedish. The teacher is only available by phone during the exam.
<b>Grades:</b>	0-13 U 14-19 3 20-25 4 26-30 5

- 1 a) Formulate Kraft-McMillan's inequality. (1 p)
- b) Explain what an instantaneous code is. (1 p)
- c) Explain how prediction with partial match (ppm) coding works. (2 p)
- d) Explain what a Golomb code is and what type of probability distribution it is good for. (2 p)
- e) Explain what the rate-distortion function is and how it is calculated for a stationary memoryless random source. (2 p)

- 2 Let  $H(X)$  be the entropy of the random variable  $X$ . Show that

$$0 \leq H(X) \leq \log L$$

where  $L$  is the size of the alphabet.

(4 p)

- 3 A memoryless source has the alphabet

$$\mathcal{A} = \{a, b, c, d, e, f, g, h\}$$

The symbol probabilities are

$$p(a) = 0.41, p(b) = 0.12, p(c) = 0.11, p(d) = 0.10$$

$$p(e) = 0.10, p(f) = 0.06, p(g) = 0.06, p(h) = 0.04$$

Construct a Huffman code for the source and calculate the resulting average data rate (in bits/symbol) of the code.

(2 p)

- 4 A stationary Markov source  $X_n$  of order 1, with alphabet  $\mathcal{A} = \{a, b, c\}$ , is given by the transition probabilities  $p(x_n|x_{n-1})$  below

$$\begin{aligned} p(a|a) &= 0.75 & p(b|a) &= 0.15 & p(c|a) &= 0.10 \\ p(a|b) &= 0.25 & p(b|b) &= 0.6 & p(c|b) &= 0.15 \\ p(a|c) &= 0.05 & p(b|c) &= 0.15 & p(c|c) &= 0.8 \end{aligned}$$

- a) Calculate the entropies  $H(X_n)$ ,  $H(X_n|X_{n-1})$  and  $H(X_n, X_{n+1}, X_{n+2})$  for the source.

(3 p)

- b) Code the sequence

*aabbcc*

using arithmetic coding. The coding should utilize the memory of the source. Give both the interval and the corresponding codeword. You can assume that the source is in state  $a$  when the coding starts and that all calculations are performed with infinite precision.

(4 p)

- 5 A source has the alphabet  $\mathcal{A} = \{m, n, o, p\}$ . A symbol sequence of length 8 is coded using BWT and mtf. The resulting index is 5 and the mtf-coded sequence is 2,0,3,0,0,1,2,0. Decode the symbol sequence.

(3 p)

6 A source has the alphabet  $\{a, b, c, d, e, f, g, h, i, j, k, l, m, n, o, p\}$ .

a) Code the sequence that begins

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using LZSS. The length of the history buffer (search buffer) is chosen as 512. Matchlengths are coded using 4 bit fixed length codewords.

(3 p)

b) A sequence from the source is coded using LZW, giving the following index sequence:

6, 0, 16, 18, 0, 12, 8, 16, 21, 6, 14, 25, 27, 4, ...

The initial dictionary is:

index	symbol	index	symbol
0	<i>a</i>	8	<i>i</i>
1	<i>b</i>	9	<i>j</i>
2	<i>c</i>	10	<i>k</i>
3	<i>d</i>	11	<i>l</i>
4	<i>e</i>	12	<i>m</i>
5	<i>f</i>	13	<i>n</i>
6	<i>g</i>	14	<i>o</i>
7	<i>h</i>	15	<i>p</i>

Decode the index sequence as far as possible. Also give the resulting dictionary.

(3 p)